

Systemgewahre statische Laufzeitanalyse von Universalbetriebssystemen

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Lehrstuhl für Informatik 4

Verteilte Systeme und Betriebssysteme

Friedrich-Alexander-Universität

Erlangen-Nürnberg

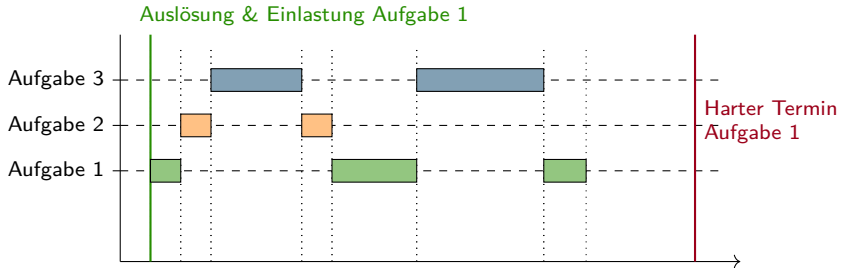


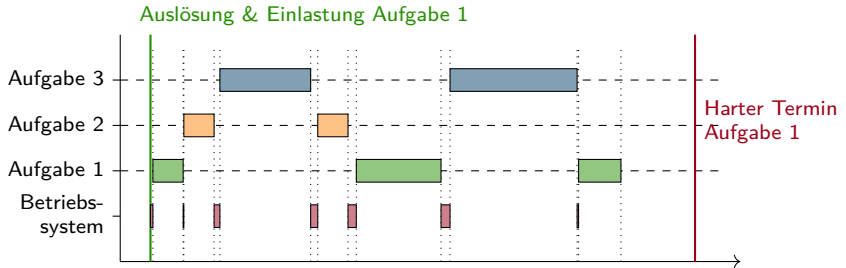
FRIEDRICH-ALEXANDER
UNIVERSITÄT
ERLANGEN-NÜRNBERG

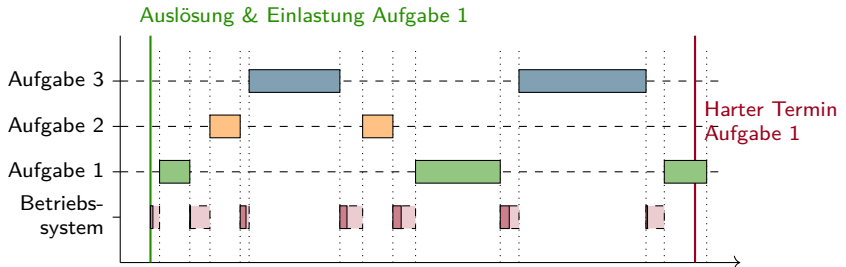


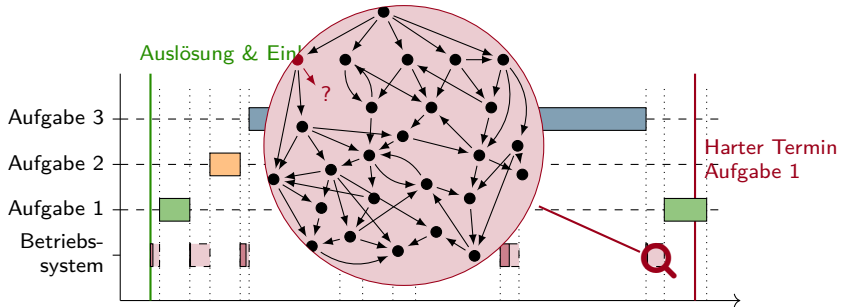
System Software Group



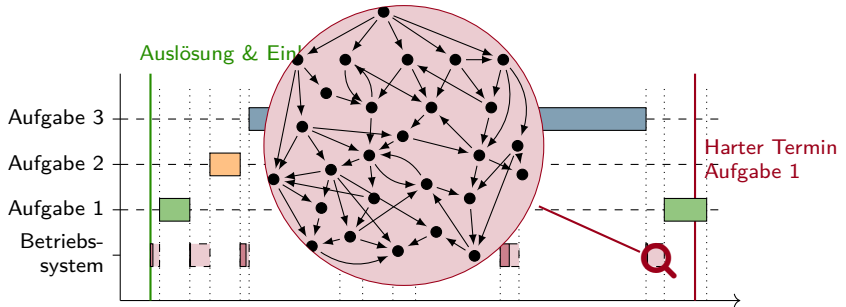




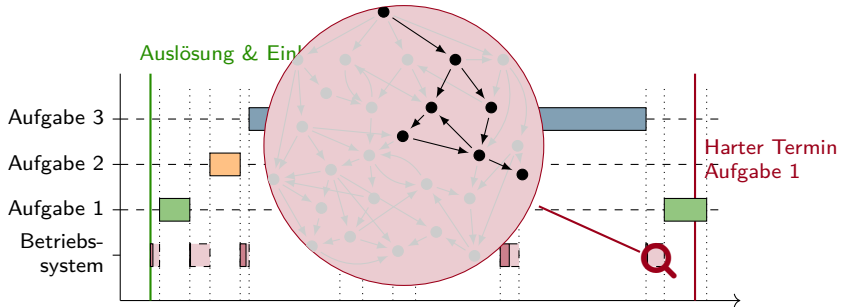




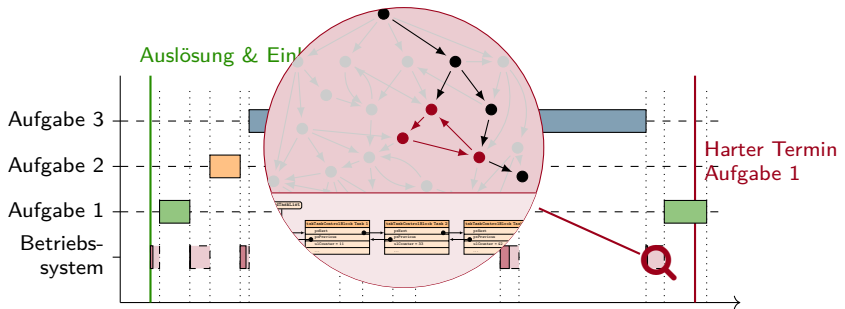
X Fehlschlagende Kontrollflussrekonstruktion durch Funktionszeiger



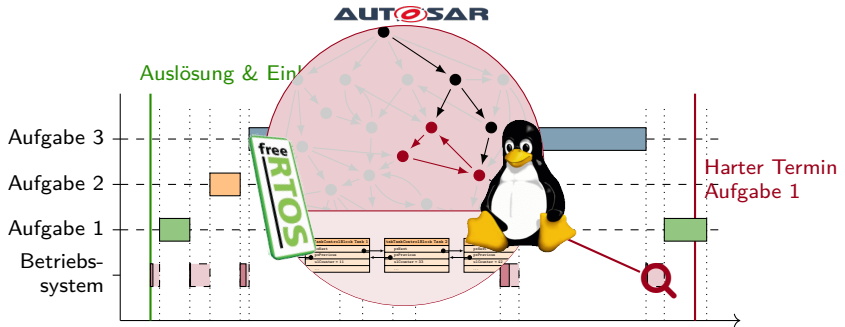
X Fehlschlagende Kontrollflussrekonstruktion durch Funktionszeiger



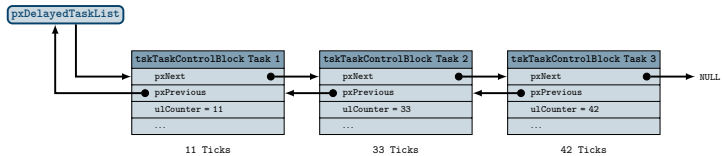
- ✗ **Fehlschlagende Kontrollflussrekonstruktion** durch Funktionszeiger
- ✗ **Hoher Pessimismus** mangels kontextsensitiver Pfadidentifikation

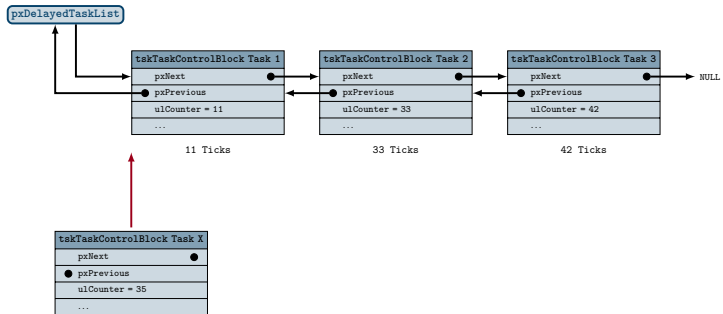


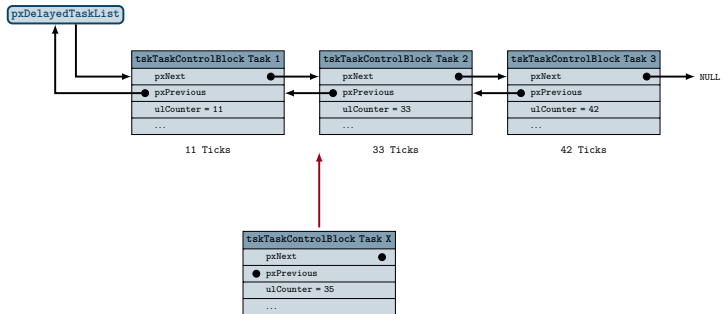
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- ✗ **Unbeschränkte WCET** aufgrund anwendungsabhängiger Datenstrukturen

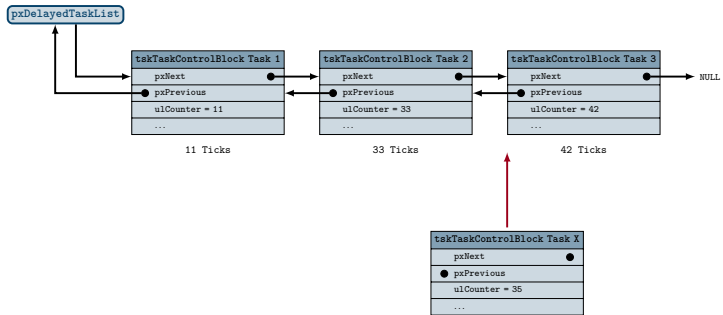


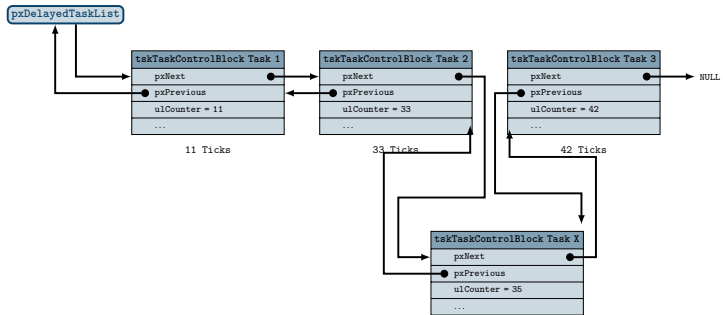
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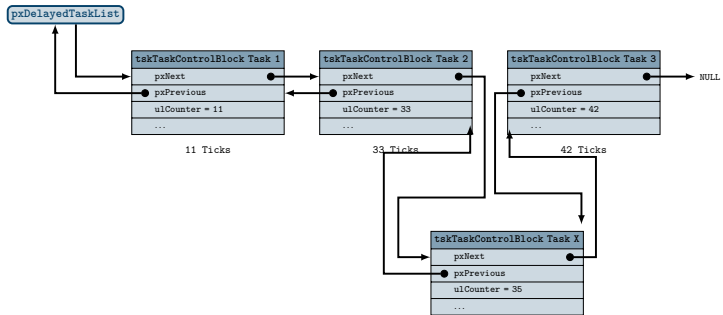


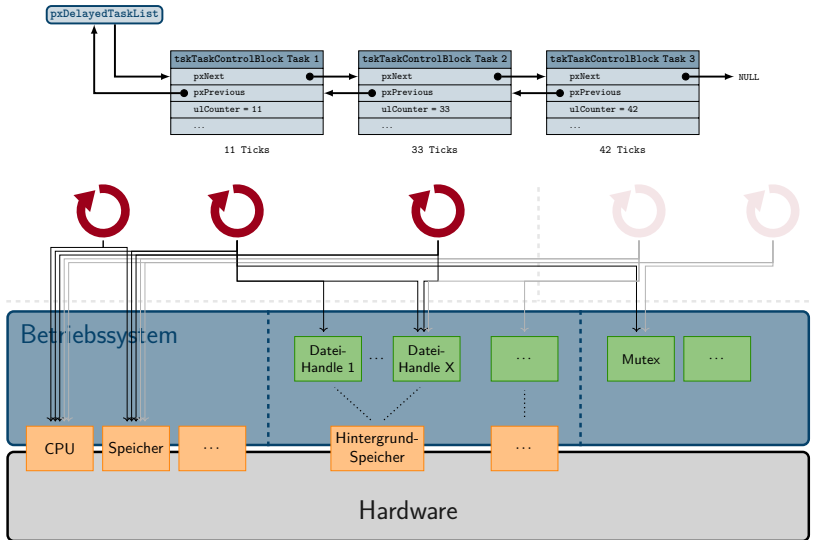


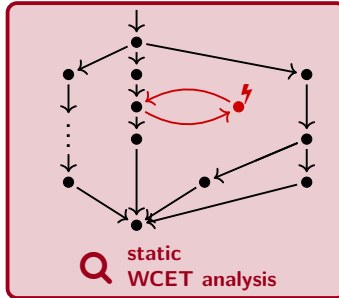


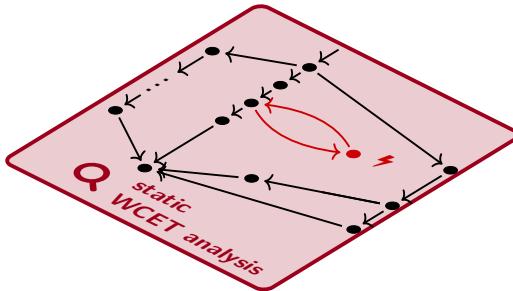


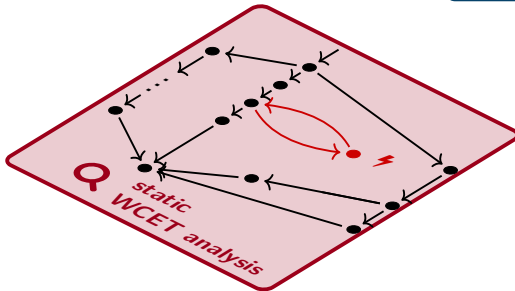
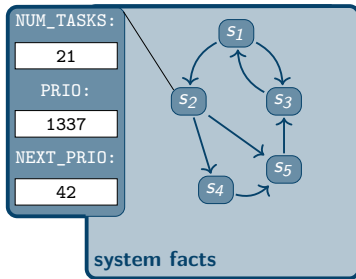




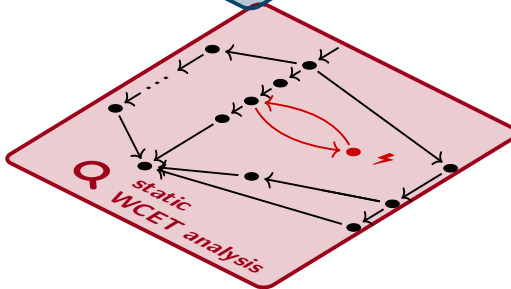
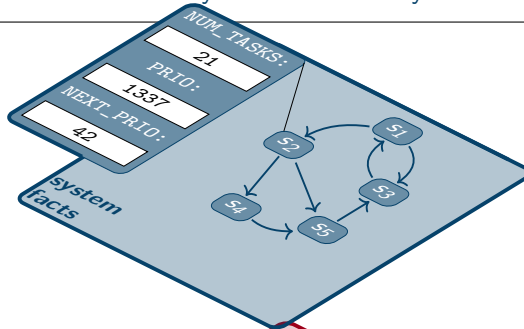


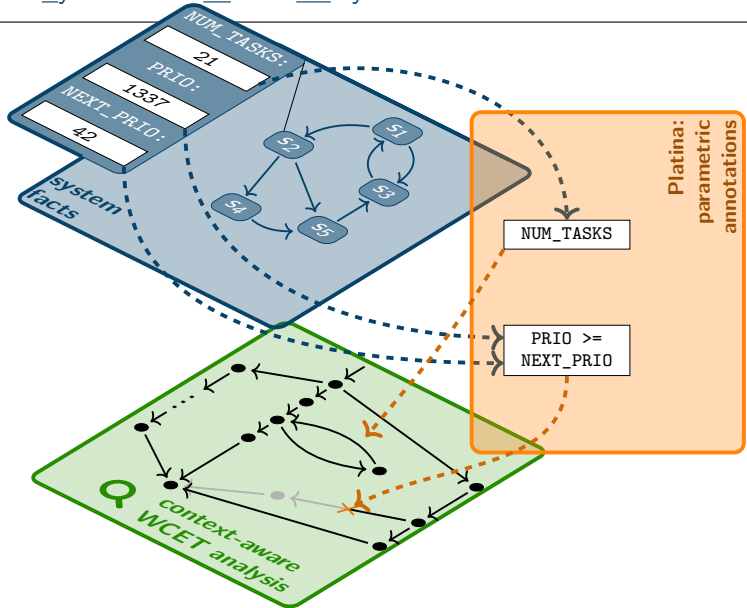


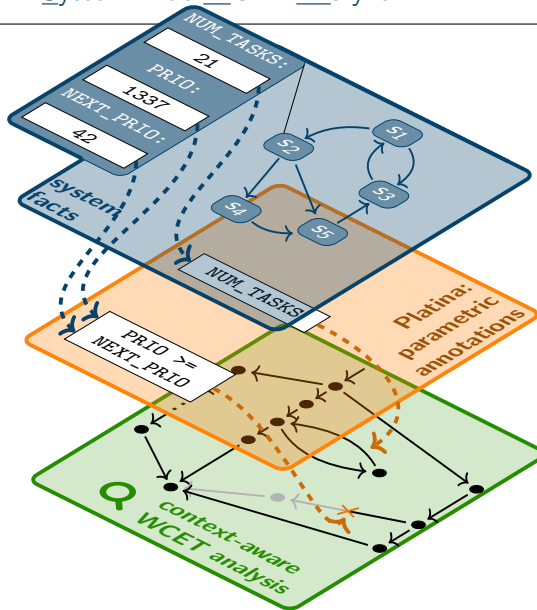


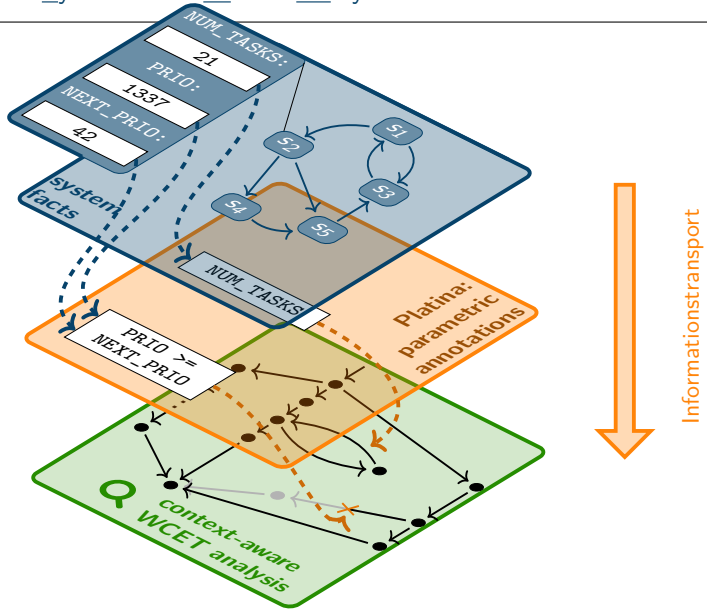


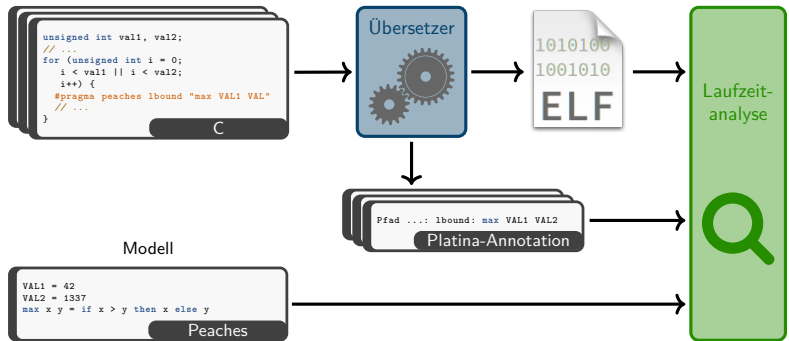
Statische Laufzeitanalyse von Betriebssystemaufrufen

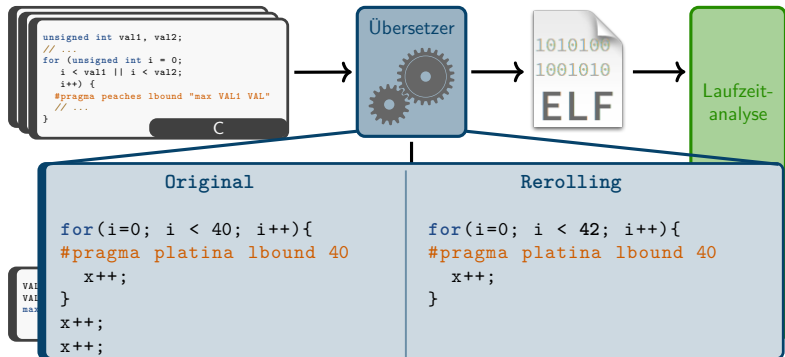


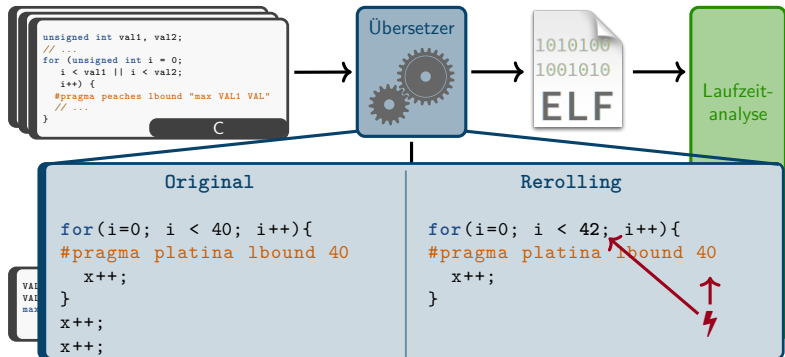


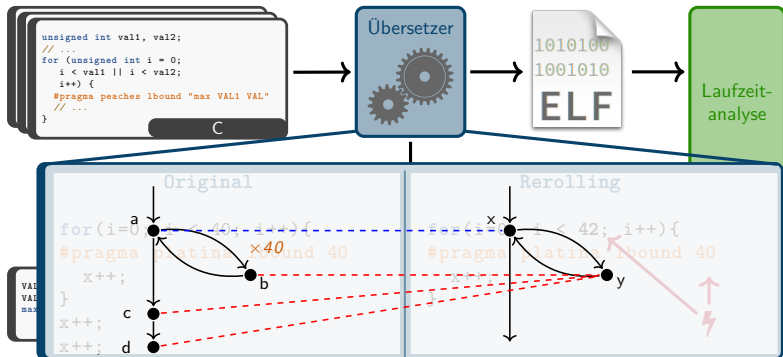


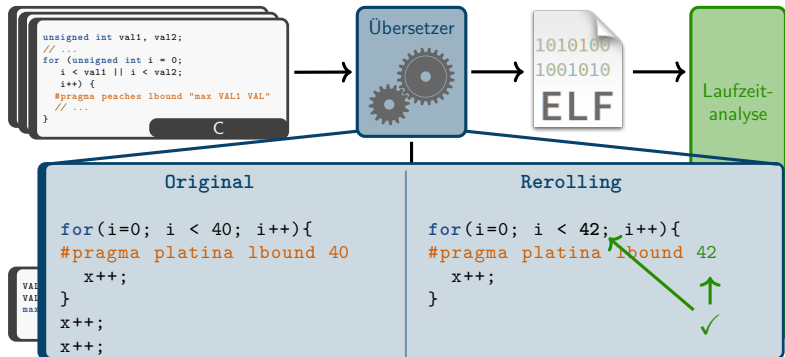


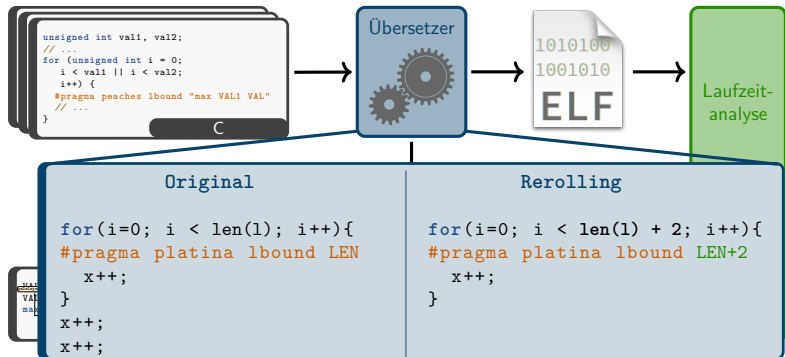


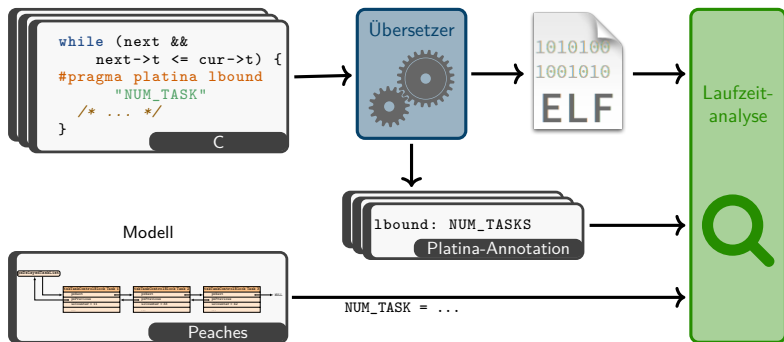


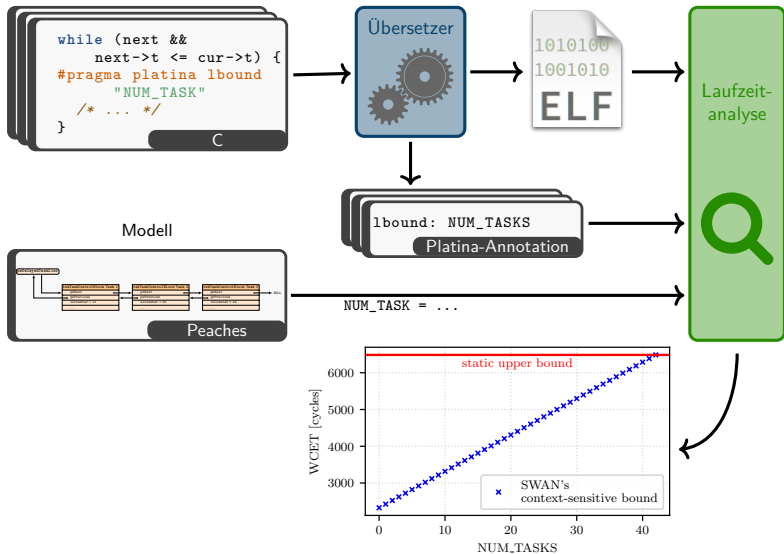


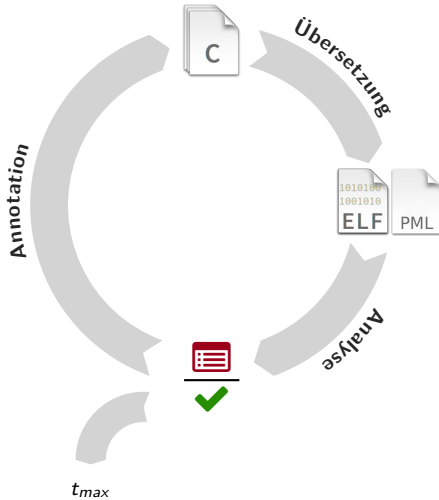


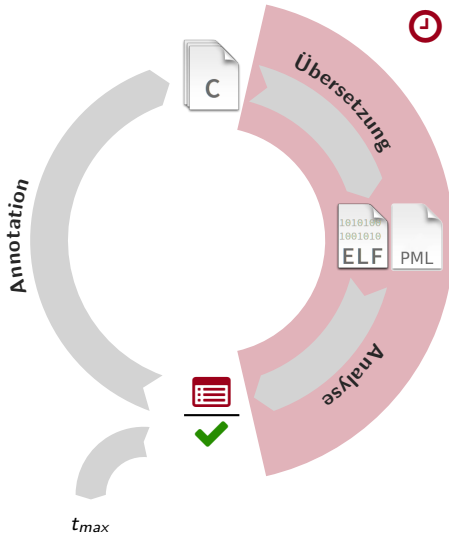






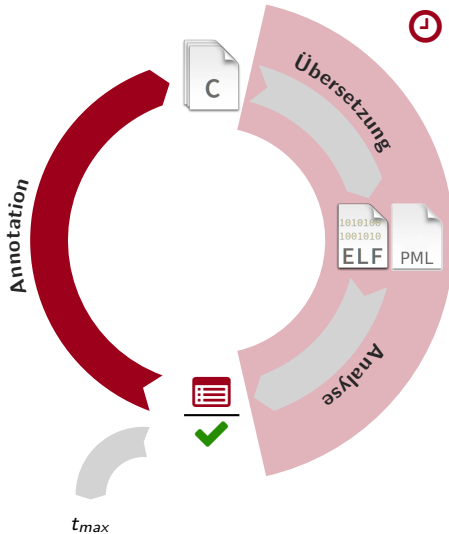






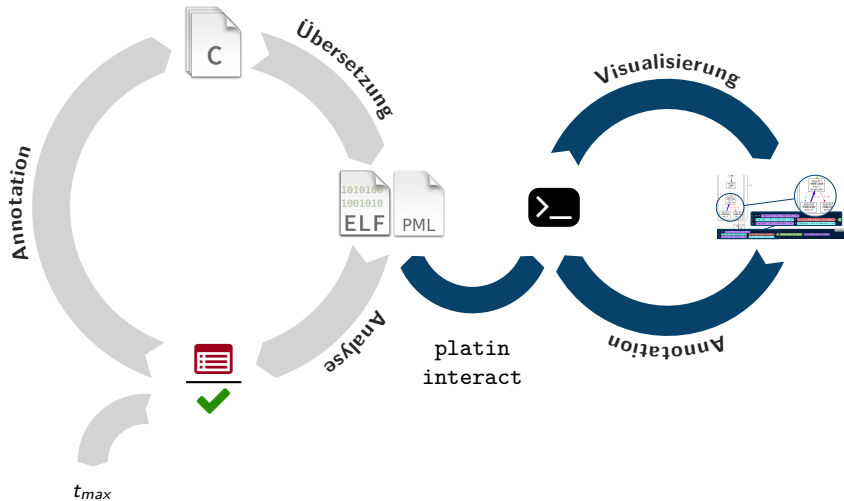
🕒 Linuxkern: über 15 min
40 Kerne @ 2.20 GHz
(Xeon E5-4640)
128GB RAM





Linuxkern: über 15 min
40 Kerne @ 2.20 GHz
(Xeon E5-4640)
128GB RAM





```
> wcet main
```



```
> wcet main
00000000
[platin] DEBUG: 74: constraint __debug_upper_bound_v20: (main)test.c:5/0/5->(loop)test.c:3 <= 10000000
[platin] DEBUG: 75: constraint __debug_upper_bound_v21: (main)test.c:5/0/9 <= 10000000
[platin] DEBUG: 76: constraint __debug_upper_bound_v22: (main)test.c:5/0/9->(guard)test.c:4 <= 10000000
[platin] DEBUG: 77: constraint __debug_upper_bound_v23: (main)test.c:5/0/13 <= 10000000
[platin] DEBUG: 78: constraint __debug_upper_bound_v24: (main)test.c:5/0/13->(callees)test.c:2 <= 10000000
[platin] DEBUG: 79: constraint __debug_upper_bound_v25: (main)test.c:5/0/17 <= 10000000
[platin] DEBUG: 80: constraint __debug_upper_bound_v26: (main)test.c:5/0/17->(faulty_annotation)test.c:0 <= 10000000
[platin] DEBUG: 81: constraint __debug_upper_bound_v27: (callees)test.c:2/0/9 <= 10000000
[platin] DEBUG: 82: constraint __debug_upper_bound_v28: (callees)test.c:2/0/9->(add)test.c:1 <= 10000000
[platin] DEBUG: #<Set:0x000000002733168>
[platin] WARNING: ILP: Unbounded loops: (guard)test.c:4/2, (loop)test.c:3/1
[platin] WARNING: Unbounded hint [(guard)test.c:4/2]: test.c:38
[platin] WARNING: Unbounded hint [(loop)test.c:3/1]: test.c:29
[platin] WARNING: WCA: ILP failed: LPSolver Error: UNBOUNDED (E3)
[platin] INFO: Finished run WCET analysis (platin) in 19 ms
[platin] INFO: best WCET bound: -1 cycles
---
- analysis-entry: main
  source: platin
  cycles: -1
>
```

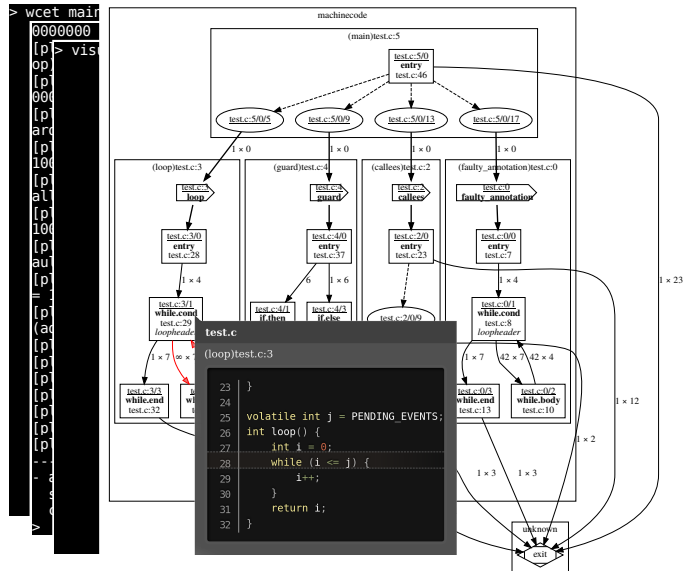


```
> wcet main
0000000
[p] visualize ilp main
```



Simon . schust

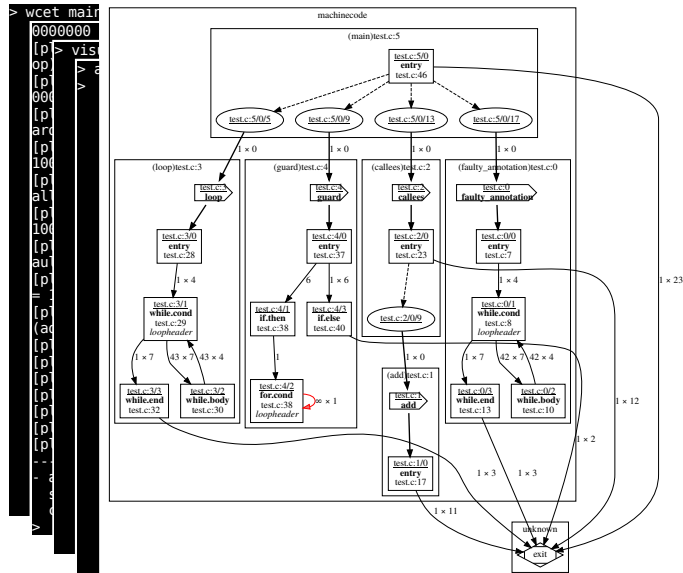





```

> wcet main
00000000
[p] > vis
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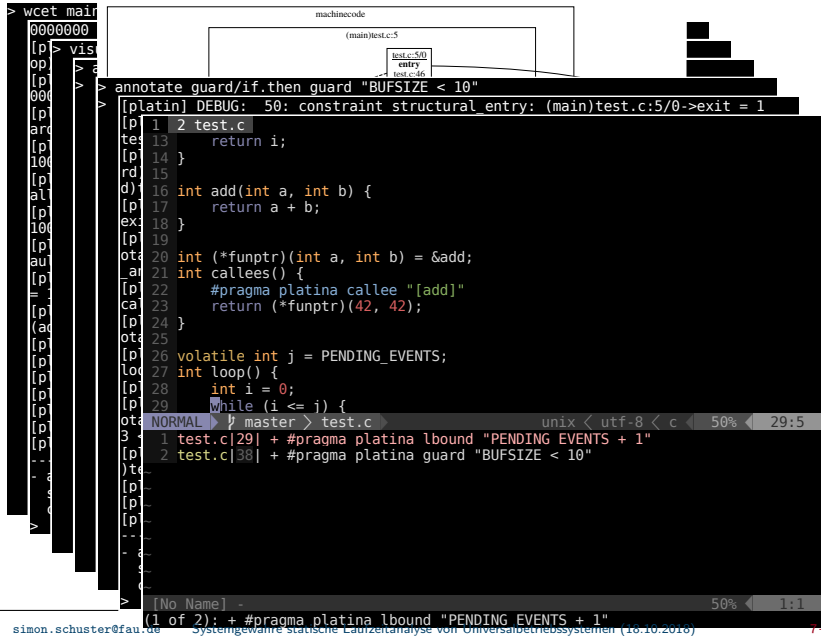
Interaktive Annotation

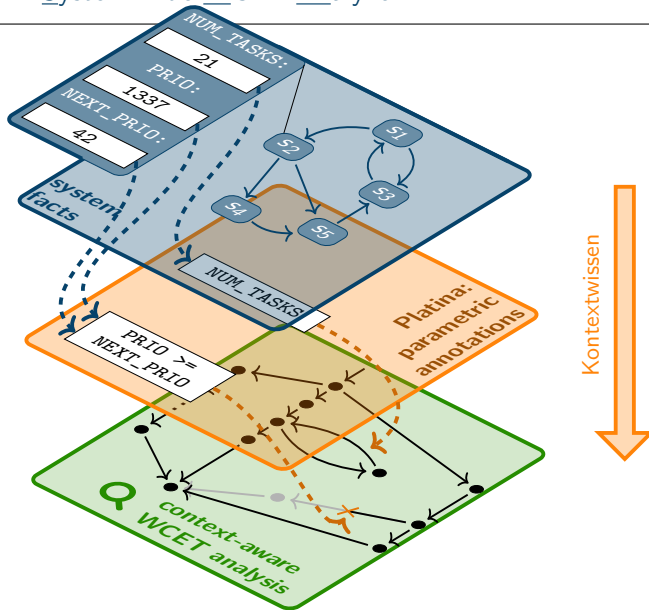
The screenshot displays a debugger's interactive annotation interface. On the left, a memory dump shows the address 00000000 and various data values. On the right, a call stack is visible with frames for 'machinecode', '(main)test.c:5', and 'test.c:5/0'. The central annotation window is titled 'annotate guard/if.then guard "BUFSIZE < 10"' and contains a large black area for editing the annotation. The interface is designed for users to interactively define and modify guards or conditions during program execution.

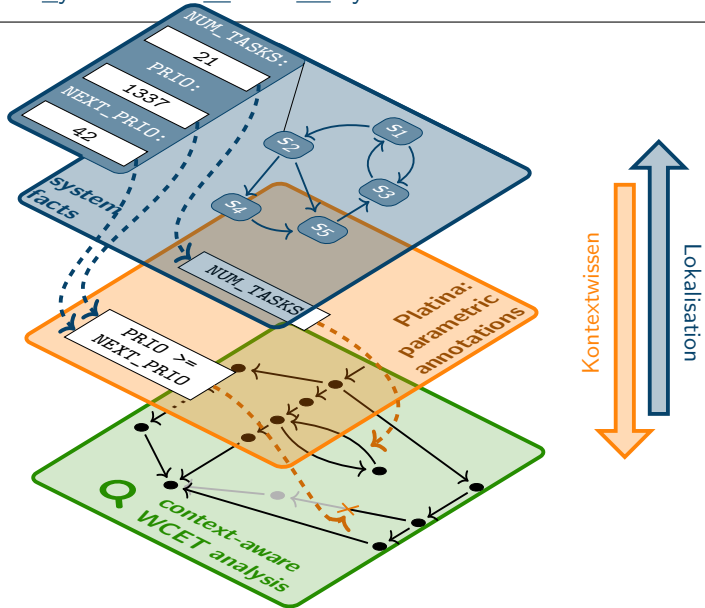




Interaktive Annotation









GatherInitiateTask (24)

```
schedlock()  
if { resume() }  
    { resume() }  
schedunlock()  
suspend(self)
```

GatherTimeoutTask (23)

```
resume()  
suspend(self)
```

GatherFinishedTask (25)

```
resume()  
resume()  
suspend(self)
```

ActuateTask (22)

```
suspend(self)
```

AttitudeTask (21)

```
suspend(self)
```




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    { resume() }  
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GatherFinishedTask (25)

```
resume()  
resume()  
suspend(self)
```

ActuateTask (22)

```
suspend(self)
```

AttitudeTask (21)

```
suspend(self)
```

```
void ResumeTask(struct tcb* next) {  
    // ...  
    if (next->prio > cur->prio) {  
        #pragma platina guard "NEXT_PRIO > CUR_PRIO"  
        reschedule(); // expensive  
    } // ...  
}
```



GatherInitiateTask (24)

```
schedlock()
if { resume()
    resume()
schedunlock()
suspend(self)
```

GatherTimeoutTask (23)

```
resume()
suspend(self)
```

GatherFinishedTask (25)

```
resume()
resume()
suspend(self)
```

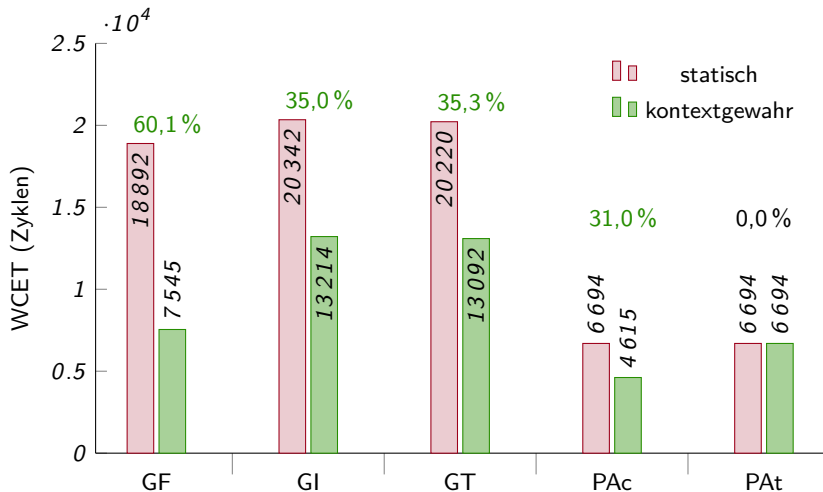
ActuateTask (22)

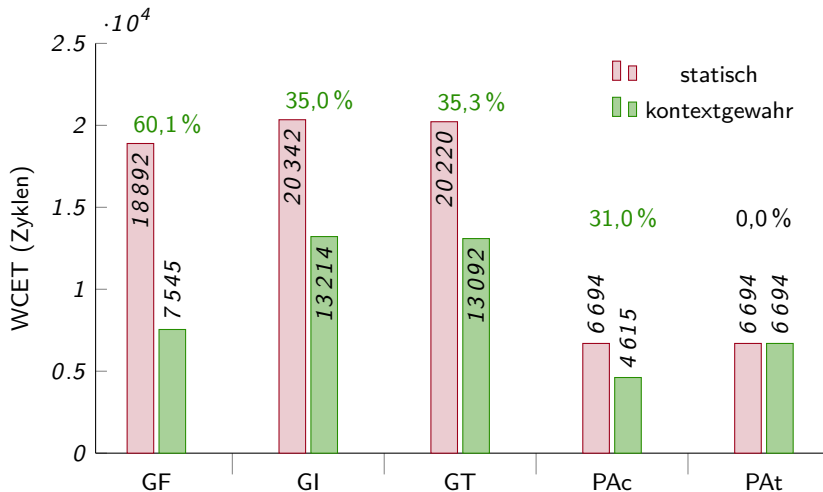
```
suspend(self)
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AttitudeTask (21)

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    } // ...
}
```





⇒ Schlimmstmögliche Antwortzeit des Aufgabensystems um **40,7 %** geringer

- statische Laufzeitanalyse von Universalbetriebssystemen problematisch
 - Pfadidentifikation
 - Abhängigkeit vom dynamischen Laufzeitkontext
- ⇒ Lösungsansatz: modellbasierte, parametrierbare statische Laufzeitanalyse
 - ✓ Annotationssprache auf Hochsprachenebene
 - ✓ Optimierungsgewahre Integration in Übersetzer und Laufzeitanalyse
 - ✓ Effiziente Annotation durch interaktive, integrierte Annotationsumgebung
- Ermöglicht statische Laufzeitanalyse kompletter Betriebssysteme
- ⇒ Kontextgewahre Analyse ermöglicht über 40 % bessere Schranken



<https://gitlab.cs.fau.de/SWAN/>