

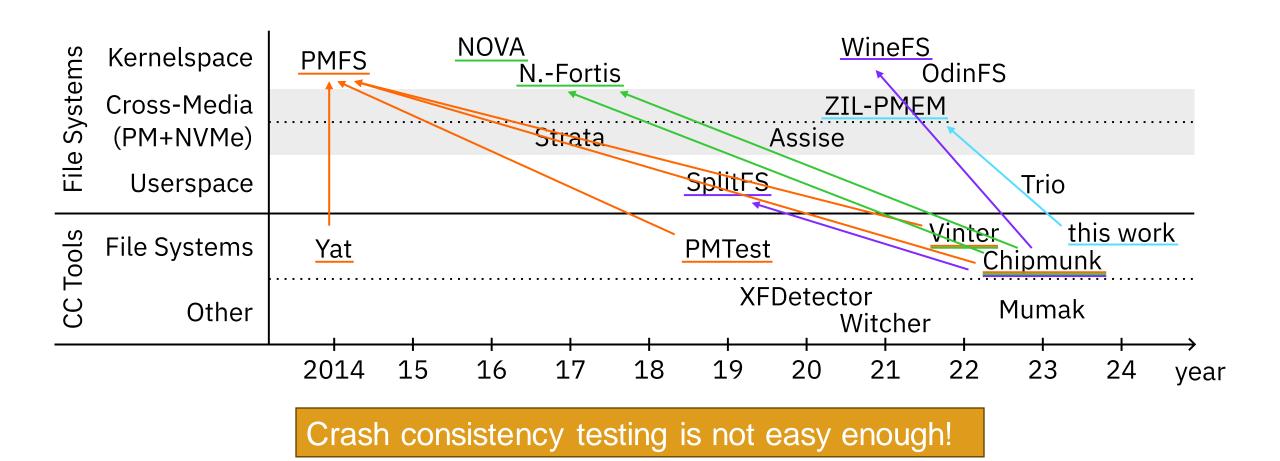
# Improvements in Crash Consistency Testing for Persistent Memory File Systems

Lukas Werling, Thomas-Christian Oder, Lucas Wäldele, Daniel Ritz, Frank Bellosa



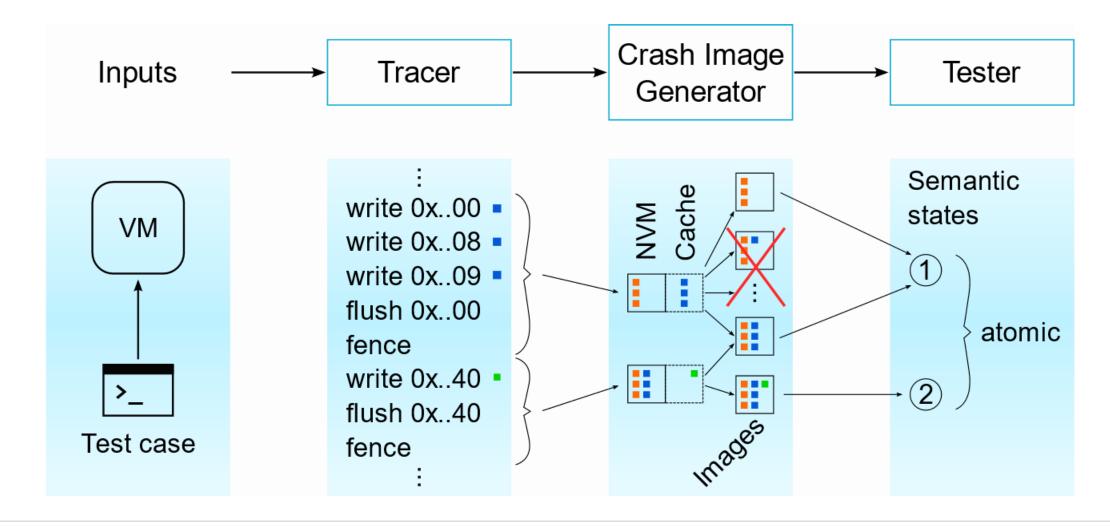
#### **Timeline**





# Background: Vinter (ATC'22)

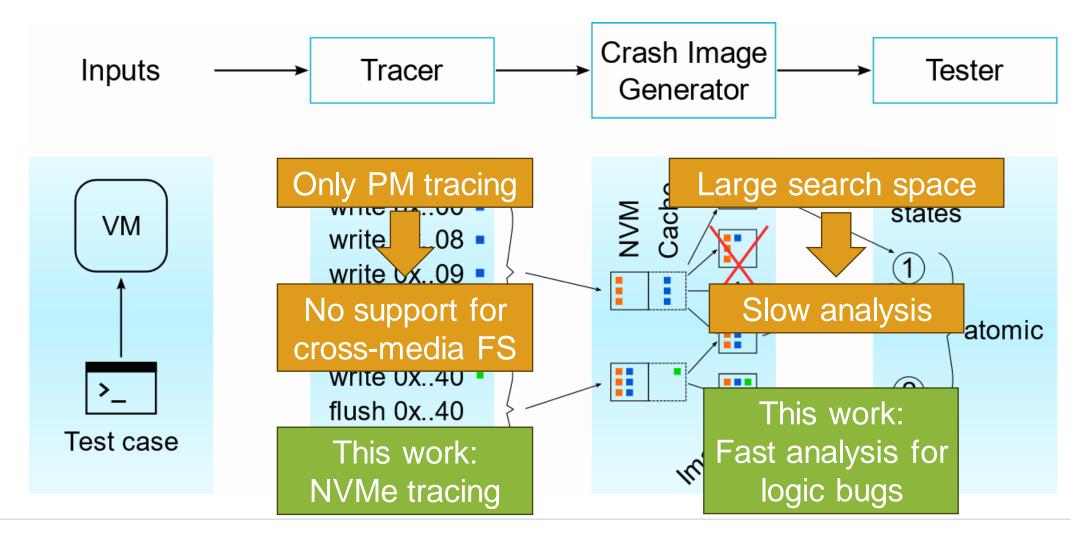




2024-03-15

## Background: Vinter (ATC'22)

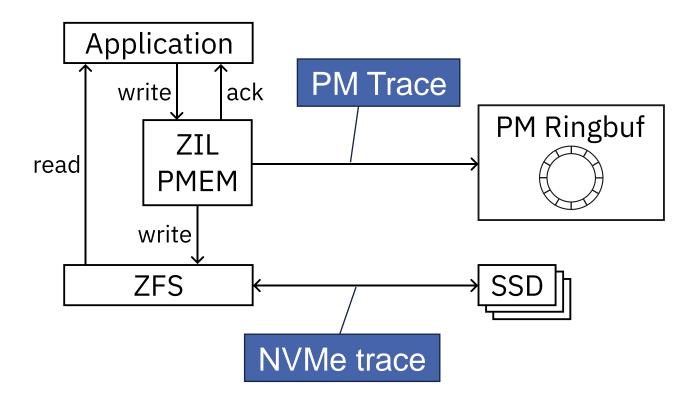




## Cross-Media File System: ZIL-PMEM

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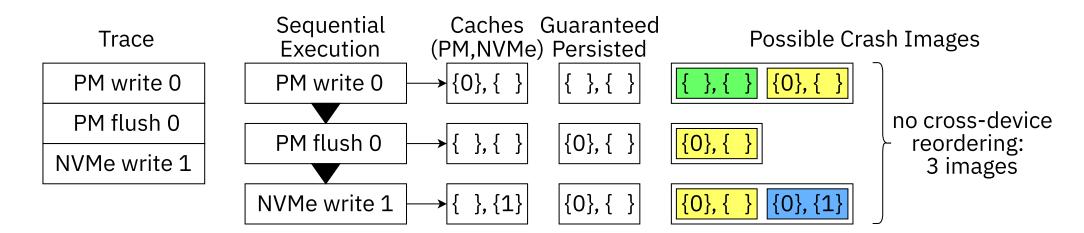




- NVMe write
  - ≈ PM stores
  - Write one or more blocks
- NVMe flush
  - ≈ PM clflush + fence
  - Commit previous writes
- NVMe is asynchronous!
  - Submission
  - Completion

## **Cross-Media Crash Images**





Is it possible to reorder PM and NVMe writes?

{ }, {1}

#### No!

- NVMe submission: UC write
- NVMe completion: interrupts
  - → memory ordering

## **Cross-Media Analysis: Implementation**



- Prototype called Permanent
- Tracer: Plain QEMU with binary translation, no PANDA
  - Hooks in virtual NVMe device
  - Hooks for memory accesses in code generation
- Crash Image Generator and Tester straightforward

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- Lots of limitations
  - No heuristic, missing optimizations, ...

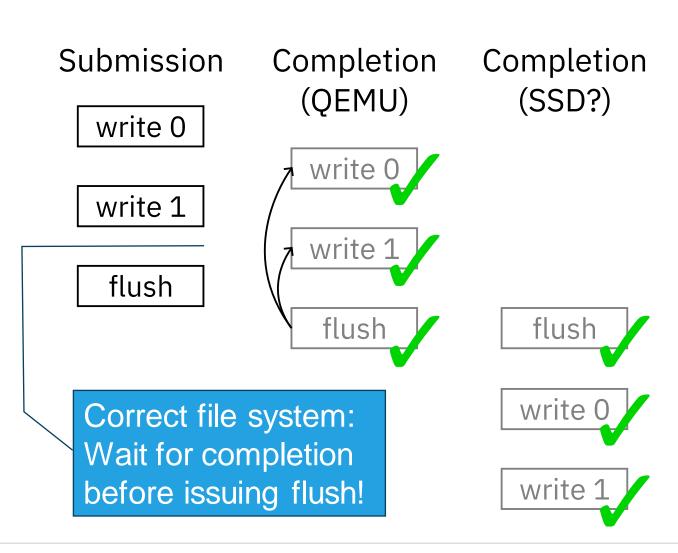
#### **Results and Discussion**



- Original Vinter file system tests
- No bugs in ZIL-PMEM!

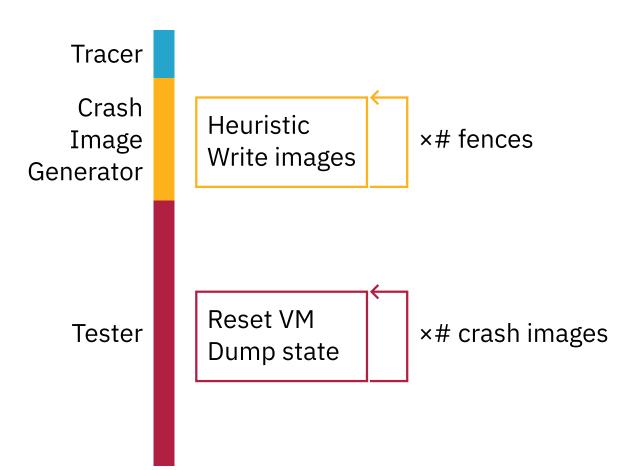
NVMe spec is not strict QEMU and real SSDs might differ

- Delayed completions
- Reordered completions



## **Vinter's Slow Analysis**





#### Goal: fewer images, similar results

- Most FS bugs are logic bugs (Chipmunk, EuroSys'23)
- Mumak (EuroSys'23)
  - Consider fewer fences
  - Generate fewer images
  - Trace analysis

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## **Vinter-FPT: Approach**



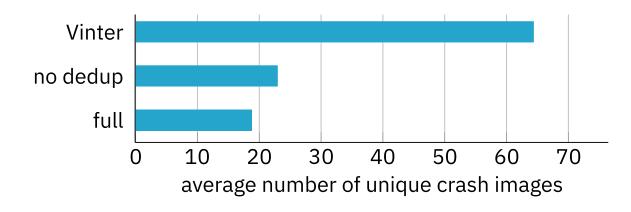
Property	Vinter	Vinter-FPT / Mumak
Where?	Memory fences	Cache flushes, memory fences
When?	Always	Once per unique call stack trace
Contents?	<ul><li>1 no stores</li><li>1 all stores</li><li>N subset of stores (heuristic)</li></ul>	1 no stores 1 all stores

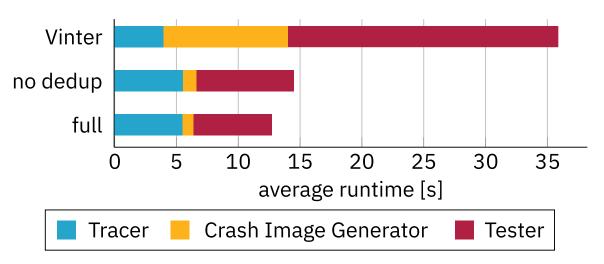
- Tracer
  - Capture stack traces
  - Small overhead!

- Crash Image Generator
  - Failure Point Tree (FPT) for deduplication
  - No more heuristic

## **Vinter-FPT: Performance**







70% fewer crash images

65% less overall runtime

Small tests suffer from slower tracer

## Vinter-FPT: Results vs. Vinter



Observation	# tests
Identical states	25
Fewer states, similar result	17
Fewer states, wrongly assumed atomic	6

Trace Analysis: 3× unordered flushes 2× missing flush

Same results for most tests

Trace analysis helps find remaining bugs

Only weak hints from trace analysis

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#### **Future Work**



#### Improve Tracing

- Binary translation costly
  - Slowdown for all instructions
  - Limited ISA support
- Idea: Selective trapping with MPK
  - Patch relevant instructions
  - Remaining code unchanged

#### **CXL Crash Consistency**

- CXL consistency model is similar
- Opportunity: Custom CXL device
- Idea: Tracing at the device
  - Observe real store ordering
  - Verify CPU consistency model

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## Conclusion



- Crash consistency testing is not as prevalent as it should be
  - Special requirements, slow analysis
- This paper: Improvements to Vinter
  - Crash consistency testing for cross-media file systems
  - Fast analysis for logic bugs

Consider crash consistency testing for your PM projects!

https://github.com/KIT-OSGroup

